A Gibbsian model for message routing in highly dense wireless networks

> András Tóbiás TU Berlin

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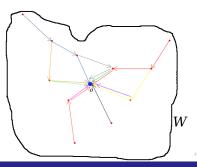
András Tóbiás (TU) A Gibbsian model for message routing

# Motivation

Consider a wireless network on a compact communication area  $W \subset \mathbb{R}^d$ . Users situated in W randomly, base station  $o \in W$ .

Idea: assume that each user sends 1 message to the base station.

- Messages travel in hops, possibly using other users as relays. Message trajectories → straight lines between consecutive steps.
- All users can take at most  $k_{\max}$  hops, for some  $k_{\max} \in \mathbb{N}$  fixed.
- A priori, message trajectories are distributed in a uniform way. All trajectories with  $1 \le k \le k_{max}$  hops are allowed, even crazy ones.



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- A priori, message trajectories are distributed in a uniform way. All trajectories with  $1 \le k \le k_{max}$  hops are allowed, even crazy ones.
- We weight this uniform distribution by 2 exponential penalty terms, preferring low interference and little congestion → Gibbsian trajectory distribution.
  - Iow interference: high signal-to-interference ratios and not too many hops,
  - little congestion: equal distribution of incoming hops among relays.
- "Common welfare" model, interplay between entropy (probability) and energy (interference+congestion).
- Question: how is the typical behaviour of trajectories (number of hops, length of a hop, shape of a trajectory) in the limit of high density of users?

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# Distribution of users: Poisson point process

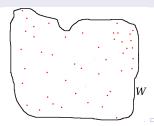
Usual assumption for wireless networks: users form a Poisson point process.

#### Definition

Let  $W \subseteq \mathbb{R}^d$  be bounded and let  $\mu$  be a finite measure on W. A random collection of points  $X = \{X_i\}_{i \in I}$  of W is a Poisson point process (PPP) in W with intensity measure  $\mu$ , if

(i)  $\forall A \subseteq W$  measurable,  $\#(X \cap A)$  is  $Poisson(\mu(A))$ distributed, i.e.,  $\mathbb{P}(\#(X \cap A) = n) = \frac{\mu(A)^n}{n!} e^{-\mu(A)}, \forall n \in \mathbb{N}_0$ ,

(ii)  $\forall k \in \mathbb{N}$ , for any pairwise disjoint sets  $A_1, \ldots, A_k \subseteq W$ , the random variables  $\{\#(X \cap A_i)\}_{i=1}^k$  are independent.



#### Communication area, users, base station

- $W \subset \mathbb{R}^d$  compact communication area, Leb(W) > 0,  $o \in W$  base station (origin of  $\mathbb{R}^d$ ).
- $\mu$  finite, absolutely continuous, nonzero measure on W.
- Users:  $X^{\lambda} = \{X_i\}_{i=1}^{N(\lambda)}$  Poisson point process with intensity  $\lambda \mu$ .
- We assume that  $(X^{\lambda})_{\lambda>0}$  is such that the empirical measure of users

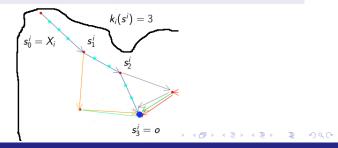
$$L_{\lambda} = \frac{1}{\lambda} \sum_{i=1}^{N(\lambda)} \delta_{X_i}$$

tends to  $\mu$  almost surely. This holds e.g. if  $\lambda \mapsto X^{\lambda}$  is increasing. (For  $x \in \mathbb{R}^d$ ,  $\delta_x$  is a measure on  $\mathbb{R}^d$ , defined via  $\delta_x(A) = 1$  if  $x \in A$  and  $\delta_x(A) = 0$  otherwise.)

### Message trajectories

- Users:  $X^{\lambda} = \{X_i\}_{i=1}^{N(\lambda)}$  Poisson point process with intensity  $\lambda \mu$ .
- Fix  $k_{\max} \in \mathbb{N}$ . Given the users  $X^{\lambda}$ , the trajectory of the message  $X_i \to o$  is random, with a random number of hops in  $\{1, \ldots, k_{\max}\}$ . It has the form

$$s^{i} = (\underbrace{k_{i}(s^{i})}_{\#\text{hops}}; \underbrace{s_{0}^{i} = X_{i}}_{\text{transmitter}}, \underbrace{s_{1}^{i} \in X^{\lambda}, \dots, s_{k_{i}(s^{i})-1}^{i} \in X^{\lambda}}_{\text{relays}}, \underbrace{s_{k_{i}(s^{i})}^{i} = o}_{\text{receiver}})$$



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#### Weighting interference

- We choose a path-loss function ℓ: [0,∞) → (0,∞): continuous, monotone decreasing, describes propagation of signal strength over distance. E.g.: ℓ(r) = min{1, r<sup>-α</sup>}, α > 0 - Hertzian propagation.
- Signal-to-interference ratio (SIR) of a transmission  $X_i \in X^{\lambda} \to x \in W$ :

$$\operatorname{SIR}_{\lambda}(X_{i}, x, X^{\lambda}) = \frac{\ell(|X_{i} - x|)}{\frac{1}{\lambda} \sum_{j=1}^{N(\lambda)} \ell(|X_{j} - x|)}$$

The denominator is called the interference at x (rescaled by  $1/\lambda$ ). • We define a SIR weight term for trajectory collections  $s = (s^i)_{i=1}^{N(\lambda)}$ :

$$\mathfrak{S}(s) = \sum_{i=1}^{N(\lambda)} \sum_{l=1}^{k_i(s')} \operatorname{SIR}_{\lambda}^{-1}(s_{l-1}^i, s_l^i, X^{\lambda}).$$

 $\rightarrow$  penalty for each step, larger if the SIR is worse (smaller).

#### Weighting congestion

- For a trajectory collection s, the number of incoming messages at the user (relay)  $X_i$  is  $m_i(s) = \sum_{i=1}^{N(\lambda)} \sum_{l=1}^{k_i(s^l)-1} \mathbb{1}\{s_l^j = X_i\}.$
- We define another weight term for the congestion:

$$\mathfrak{M}(s) = \sum_{i=1}^{N(\lambda)} m_i(s)(m_i(s)-1).$$

 $\rightarrow$  number of ordered pairs of incoming messages at all relays. Large penalty for uneven distributions of incoming messages among relays.

# The Gibbs distribution

#### Definition of the Gibbs distribution

For the intensity  $\lambda > 0$  and two parameters  $\gamma > 0$ ,  $\beta \ge 0$ , given the users  $X^{\lambda} = (X_i)_{i=1}^{W(\lambda)}$ , the message trajectories are chosen according to the following Gibbs distribution:

$$\mathrm{P}_{\lambda,X^{\lambda}}^{\gamma,\beta}(s) = \frac{1}{Z_{\lambda}^{\gamma,\beta}(X^{\lambda})} \frac{1}{N(\lambda)^{\sum_{i=1}^{N(\lambda)}(k_i(s^i)-1)}} \exp(-\gamma \mathfrak{S}(s) - \beta \mathfrak{M}(s))$$

Here  $Z_{\lambda}^{\gamma,\beta}(X^{\lambda})$  is the normalizing constant, called partition function, which makes  $P_{\lambda,X^{\lambda}}^{\gamma,\beta}$  a probability measure:

$$Z_{\lambda}^{\gamma,\beta}(X^{\lambda}) = \sum_{r} \frac{1}{N(\lambda) \sum_{i=1}^{N(\lambda)} (k_i(r^i) - 1)} \exp(-\gamma \mathfrak{S}(r) - \beta \mathfrak{M}(r)).$$

#### Plan to analyze the high-density limit $\lambda \to \infty$

Given  $(X^{\lambda})_{\lambda>0}$ , determine the limiting free energy  $\lim_{\lambda\to\infty} \frac{1}{\lambda} \log Z_{\lambda}^{\gamma,\beta}(X^{\lambda})$ . The free energy is expected to be given by a variational formula  $\rightarrow$ minimizer(s) give information about the limiting distribution of trajectories.

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### First approach to the limiting free energy

Idea: use the empirical measures of message trajectories of given lengths k. For  $k = 1, ..., k_{\max}$  and for a trajectory collection  $s = (s^i)_{i=1}^{N(\lambda)}$ , we put $R_{\lambda,k}(s) = \frac{1}{\lambda} \sum_{i=1}^{N(\lambda)} \delta_{(s_0^i,...,s_{k-1}^i)} \mathbb{1}\{k_i(s^i) = k\}.$ 

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#### Properties

- For k = 1,..., k<sub>max</sub> and for all s, R<sub>λ,k</sub>(s) is a random element of the set M(W<sup>k</sup>) of finite measures on W<sup>k</sup> = W<sup>{0,1,...,k-1}</sup>.
- The partition function  $Z_{\lambda}^{\gamma,\beta}(X^{\lambda})$  is a function of these measures.
- Each user sends 1 message to  $o \Rightarrow$  the 0th marginals  $\pi_0 R_{\lambda,k}(s)$  of the  $R_{\lambda,k}(s)$ 's sum up to the empirical measure of users  $L_{\lambda}$ :

$$\sum_{k=1}^{k_{\max}} \pi_0 R_{\lambda,k}(s) = rac{1}{\lambda} \sum_{i=1}^{N(\lambda)} \delta_{X_i} = L_\lambda.$$

- Assumed:  $L_{\lambda} \Rightarrow \mu$ , almost surely. Thus, along a subsequence, the  $R_{\lambda,k}(\cdot)$ 's converge to some  $\Sigma = (\nu_k)_{k=1}^{k_{\text{max}}}$ ,  $\nu_k \in \mathcal{M}(W^k)$ , with  $\sum_{k=1}^{k_{\text{max}}} \pi_0 \nu_k = \mu$ .
- The SIR term  $\mathfrak{S}(\cdot)$  depends on each  $R_{\lambda,k}(\cdot)$  in a continuous, linear way.

# The congestion term makes trouble

Problem: the congestion term depends discontinuously on  $R_{\lambda,k}(\cdot)$ 's as  $\lambda \to \infty$ .  $\longrightarrow$  No way to express its limit in the terms of the limiting measures  $(\nu_k)_{k=1}^{k_{max}}$ .

### The congestion term makes trouble

**Problem:** the congestion term depends discontinuously on  $R_{\lambda,k}(\cdot)$ 's as  $\lambda \to \infty$ .  $\longrightarrow$  No way to express its limit in the terms of the limiting measures  $(\nu_k)_{k=1}^{k_{max}}$ . Solution: we introduce the empirical measures of users receiving given numbers of incoming messages w.r.t. the trajectory family *s*:

$$P_{\lambda,m}(s) = \frac{1}{\lambda} \sum_{i=1}^{N(\lambda)} \delta_{X_i} \mathbb{1}\{m_i(s) = m\}, \quad m \in \mathbb{N}_0.$$
(1)

#### Properties

• Each  $P_{\lambda,m}(s)$  is a random element of  $\mathcal{M}(W)$ .

• The congestion term  $\mathfrak{M}(\cdot)$  depends linearly on this family of measures:

$$\mathfrak{M}(s) = \sum_{i=1}^{N(\lambda)} m_i(s)(m_i(s)-1) = \lambda \sum_{m=0}^{\infty} m(m-1)P_{\lambda,m}(s)(W).$$

- Since each user receives exactly *m* incoming messages for precisely one *m*,  $\sum_{m=0}^{\infty} P_{\lambda,m}(s) = L_{\lambda}, \quad \forall s, \lambda.$
- So  $(P_{\lambda,m}(\cdot))_m$  also converge along a subsequence to some  $\Xi = (\mu_m)_{m=0}^{\infty}$ , where  $\mu_m \in \mathcal{M}(W)$ , with  $\sum_{m=0}^{\infty} \mu_m = \mu$ .

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Form of limiting measures:  $\Psi = (\Sigma, \Xi) = ((\nu_k)_{k=1}^{k_{\max}}, (\mu_m)_{m \in \mathbb{N}_0})$ :

- $\nu_k \in \mathcal{M}(W^k)$ ,  $k = 1, \ldots, k_{max}$ : limiting distribution of k-hop trajectories,
- $\mu_m \in \mathcal{M}(W)$ ,  $n \in \mathbb{N}_0$ : limiting distribution of users (relays) receiving precisely *m* incoming messages,

#### Constraints

- (i)  $\sum_{k=1}^{k_{\text{max}}} \pi_0 \nu_k = \mu$  because each user sends out 1 message to o,
- (ii)  $\sum_{m=0}^{\infty} \mu_m = \mu$  because each user receives *m* incoming messages for exactly one *m*,
- (iii)  $\sum_{k=1}^{k_{\text{max}}} \sum_{l=1}^{k-1} \pi_l \nu_k = \sum_{m=0}^{\infty} m \mu_m$ : the total number of relaying hops of all trajectories = the total number of incoming messages at all relays.

# The limiting free energy

#### Theorem

We have for  $\beta \geq 0$ ,  $\gamma > 0$ , almost surely w.r.t. the users  $(X^{\lambda})_{\lambda > 0}$ ,

$$\lim_{\lambda \to \infty} \frac{1}{\lambda} \log Z_{\lambda}^{\gamma,\beta}(X^{\lambda}) = - \inf_{\Psi \text{ satisfying (i), (ii), (iii)}} (\mathrm{I}(\Psi) + \gamma \mathrm{S}(\Psi) + \beta \mathrm{M}(\Psi)).$$

• 
$$\Psi = (\Sigma, \Xi) = ((\nu_k)_{k=1}^{k_{\max}}, (\mu_m)_{m=0}^{\infty})$$
 satisfying (i)–(iii).

- $S(\Psi) = S(\Sigma)$ : a limiting SIR term depending only on the  $\nu_k$ 's.
- $M(\Psi) = M(\Xi)$ : a limiting congestion term depending only on the  $\mu_m$ 's.
- I(Ψ): an entropy term → logarithmic rate of combinatorial terms expressing counting complexity. Involves both Σ and Ξ.
- (Precise expressions for S, M, I are on the last slide.)
- We'll see: the variational formula has at least 1 minimizer.

# Analysis of the minimizers

Strategy: show that a minimizer exists + all minimizers are positive wherever  $\mu$  is positive  $\rightarrow$  identify minimizers via deriving the Euler–Lagrange equations.

Case  $\beta > 0$ ,  $\gamma > 0$ 

Uniqueness is unclear. All minimizers are given in the following implicit way: for  $x, x_0, \ldots, x_{k-1} \in W$ ,

$$\begin{split} \nu_k(\mathrm{d} x_0, \dots, \mathrm{d} x_{k-1}) &= \mu(\mathrm{d} x_0) A(x_0) \prod_{l=1}^{k-1} C(x_l) M(\mathrm{d} x_l) \mathrm{e}^{-\gamma \frac{\int_W \ell(|z-x_l|)\mu(\mathrm{d} x_l)}{\ell(|x_l-1-x_l|)}},\\ \mu_m(\mathrm{d} x) &= \mu(\mathrm{d} x) B(x) \frac{(C(x)\mu(W))^{-m}}{m!} \mathrm{e}^{-\beta m(m-1)}. \end{split}$$

Here A, B, C are positive functions s.t. (i),(ii),(iii) hold,  $M = \sum_{k=1}^{k_{max}} \sum_{l=1}^{k-1} \pi_l \nu_k = \sum_{m=0}^{\infty} m \mu_m.$ 

#### Case $\beta = 0$ , $\gamma > 0$

These equations remain true, but they simplify + uniqueness holds. Can write

$$\nu_k(\mathrm{d} x_0,\ldots,\mathrm{d} x_{k-1}) = \mu(\mathrm{d} x_0) A(x_0) \prod_{l=1}^{k-1} \frac{\mu(\mathrm{d} x_l)}{\mu(W)} \mathrm{e}^{-\gamma \frac{\int_W \ell(|z-x_l|) \mu(\mathrm{d} x_l)}{\ell(|x_{l-1}-x_l|)}}.$$

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#### Law of large numbers for the empirical measures of trajectories:

if  $\gamma > 0$  and  $\beta = 0$ , the empirical measures of trajectories  $(R_{\lambda,k}((S^i)_{i=1}^{N(\lambda)})_{k=1}^{k_{max}}$  converge to the unique minimizer  $(\nu_k)_{k=1}^{k_{max}}$  of the variational formula. This follows from a large deviation principle for these empirical measures.

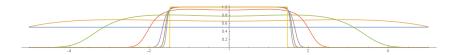
The minimizer is amenable for analytical investigations  $\rightarrow$  gives information about the network for high user densities  $\lambda < \infty$ .

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 $\nu_k$ : limiting distribution of k-hop trajectories.

# Qualitative properties of the network

Law of large numbers for the empirical measures of trajectories: if  $\gamma > 0$  and  $\beta = 0$ , the empirical measures of trajectories  $(R_{\lambda,k}((S^i)_{i=1}^{N(\lambda)})_{k=1}^{k_{\text{max}}}$  converge to the unique minimizer  $(\nu_k)_{k=1}^{k_{\text{max}}}$  of the variational formula.  $\nu_k$ : limiting distribution of k-hop trajectories.



#### Example: one-hop trajectories in a one-dimensional setting

Density of 1-hop trajectories  $\nu_1(dx)/\mu(dx)$  for  $\gamma = 0, 0.001, 0.01, 0.1, 1, \infty$ , for  $W = [-5,5] \subset \mathbb{R}$ ,  $\mu = \text{Leb}|_W$ , o = 0,  $\ell(r) = \min\{1, r^{-4}\}$ ,  $k_{\max} = 2$ . For  $\gamma$  close to 0,  $\nu_1$  is almost identically 1/2. For  $\gamma$  large enough (already for  $\gamma = 1$ !),  $\nu_1(dx_0)/\mu(dx_0)$  is close to the indicator function of the 1-hop path being better w.r.t. SIR penalization than any of the 2-hop paths from  $x_0$  to o.

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#### 1. Typical number of hops in a large-distance limit

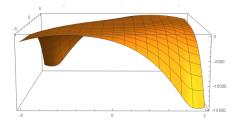
Consider larger and larger balls W with users distributed according to Lebesgue measure  $\mu$ ,  $k_{max}$  large, transmitter  $x_0$  far out. Under suitable assumptions on  $\ell$ , the typical length of a hop tends to infinity! E.g., if  $\ell(r) = \min\{1, r^{-\alpha}\}, \alpha > d$ , then  $x_0 \to o$  typically takes  $k(|x_0|) \asymp \frac{|x_0|}{\log^{1/\alpha} |x_0|}$  hops of equal length  $\asymp \log^{1/\alpha} |x_0|$ . The optimal path follows a straight line with equal-sized hops, macroscopic deviations from it get exponentially unlikely in this limit. Law of large numbers for the empirical measures of trajectories: if  $\gamma > 0$  and  $\beta = 0$ , the empirical measures of trajectories  $(R_{\lambda,k}((S^i)_{i=1}^{N(\lambda)})_{k=1}^{k_{max}}$  converge to the unique minimizer  $(\nu_k)_{k=1}^{k_{max}}$  of the variational formula.  $\nu_k$ : limiting distribution of k-hop trajectories.

#### 2. Convergence to the straight line for fixed W and large $\gamma$

Fix  $k_{\max}$  and the communication area  $W = \overline{B_r(o)}$ , let  $\mu$  be rotationally symmetric and  $\ell$  strictly monotone increasing. E.g.:  $\ell(r) = (1 + r)^{-\alpha}$ . Then as  $\gamma \to \infty$ , we observe convergence to the straight line: for any  $\varepsilon > 0$ ,  $\forall x_0 \in W$ , the probability of choosing trajectories  $x_0 \to o$  with  $\geq 1$  hop  $\geq \varepsilon$  away from the straight line decays exponentially fast.

# Thank you for your attention!

- W. König and A. Tóbiás: A Gibbsian model for highly dense multihop networks. arXiv:1704.03499 (2017) – for the general case (penalizing interference+congestion).
- W. König and A. Tóbiás: Routeing properties in a Gibbsian model for highly dense multihop networks. arXiv:1801.04985 (2017/18) – for the applications: qualitative properties of the network, motivation, game-theoretic properties, simulation results.



### The limiting entropy, SIR and congestion terms

For  $\Psi = ((\nu_k)_{k=1}^{k_{max}}, (\mu_m)_{m=0}^{\infty})$ , the entropy term  $I(\Psi)$  is given as

$$\begin{split} \mathrm{I}(\Psi) &= \sum_{k=1}^{k_{\max}} \int_{W^k} \nu_k(\mathrm{d}x_0, \dots, \mathrm{d}x_{k-1}) \log \frac{\mathrm{d}\nu_k}{\mathrm{d}\mu^{\otimes k}}(x_0, \dots, x_{k-1}) \\ &+ \sum_{m=0}^{\infty} \int_{W} \mu_m(\mathrm{d}x) \log \frac{\mathrm{d}\mu_m}{\mathrm{d}\mu c_m}(x) - \int_{W} M(\mathrm{d}x) \log \frac{\mathrm{d}M}{\mathrm{d}\mu}(x) - \frac{1}{\mathrm{e}}, \end{split}$$

where  $c_m$  are the weights of a Poisson $\left(\frac{1}{e_u(W)}\right)$ -distribution.

The expression is to be understood as  $+\infty$  if some of the Radon-Nikodym derivatives doesn't exist, and we use the convention  $0 \log 0 = 0 \log \frac{0}{0} = 0$ . The limiting SIR term is

$$\mathrm{S}(\Psi) = \sum_{k=1}^{k_{\max}} \sum_{l=1}^{k} \int_{W^k} \nu_k(\mathrm{d} x_0, \ldots, \mathrm{d} x_{k-1}) \frac{\int_W \ell(|z - x_l|) \mu(\mathrm{d} z)}{\ell(|x_{l-1} - x_l|)} \in [0, \infty).$$

The limiting congestion term is

$$\mathrm{M}(\Psi) = \sum_{m=0}^{\infty} m(m-1)\mu_m(W) \in [0,\infty].$$

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